

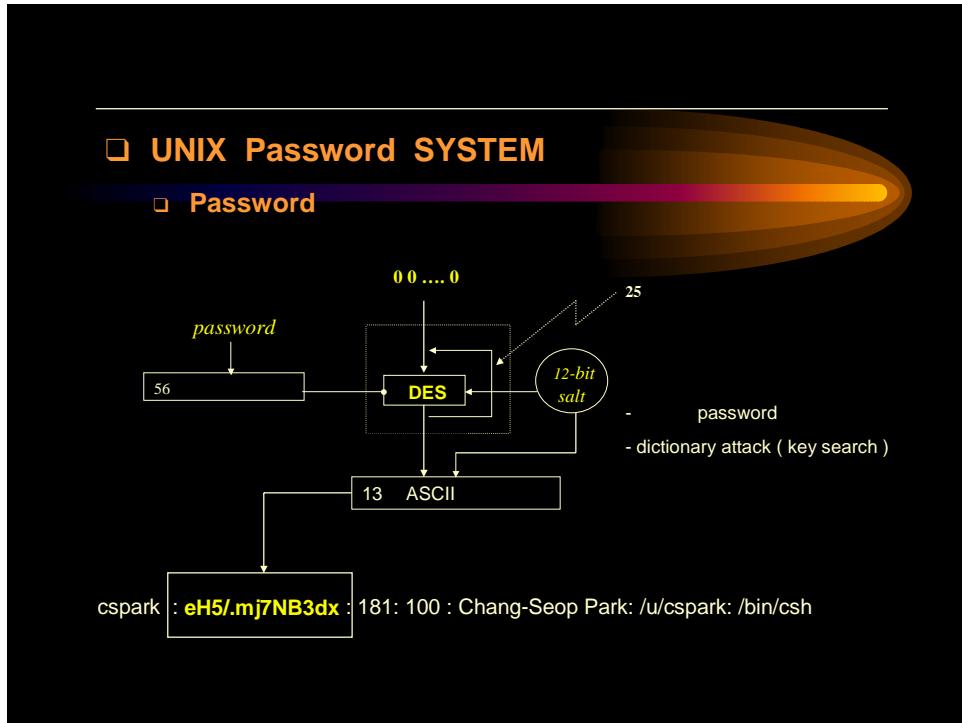
□ **Password**

□ **Password Control**

- Limited Attempts ()
- Password Aging ()
- Minimum Length ()
- Lockout of Dormant Account (account)

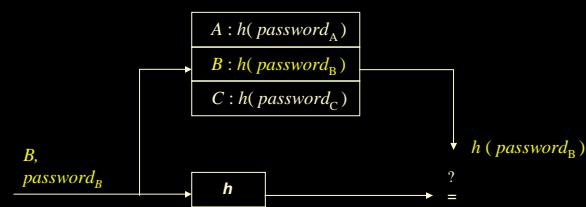
□ **Password Checking for Good Password**

- Reactive Password Checking - "CRACK"
- Proactive Password Checking

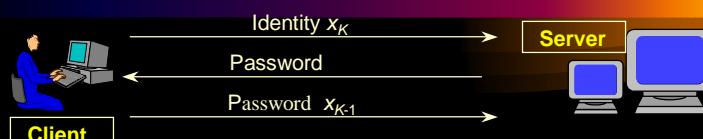


❑ UNIX Password SYSTEM

❑ Password

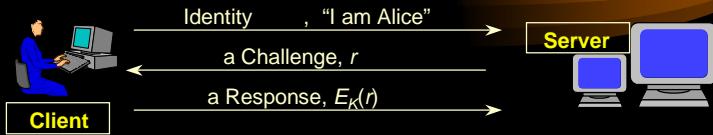


❑ One-Time Password



- ❑ System Password
- ❑ Client Server Secret
- ❑ **One-way Function $h()$**
- ❑ Client $x_0 \quad x_i = h(x_{i-1})$ for $i = 1, 2, 3, \dots, k$ { k access tickets }
- ❑ password x_k access user identity
- ❑ System
- ❑ Sequence x_i System

□ Challenge-Response



- -
- Client Server가 secret key K
- E is a public encryption function e.g. DES ; hash function
- (secret information)
- key database
- Response ; E

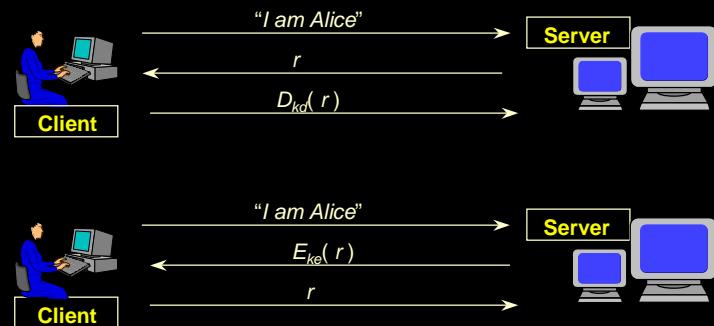
□ Challenge-Response

□ Symmetric Key

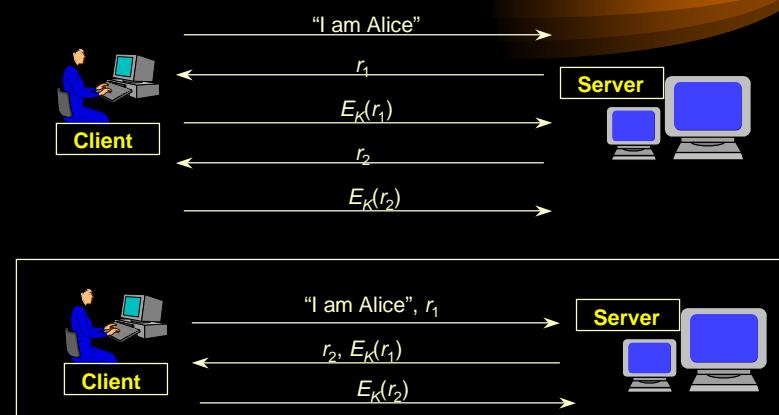


□ Challenge-Response

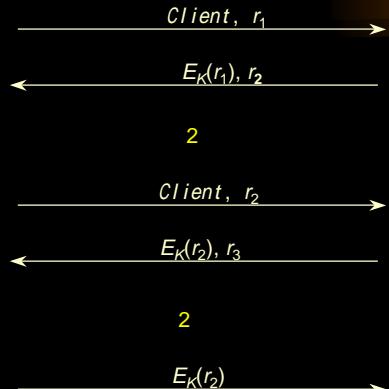
□ Public Key



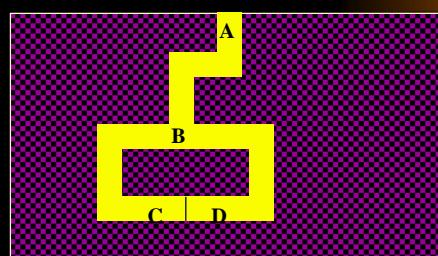
□ Mutual Identification ()



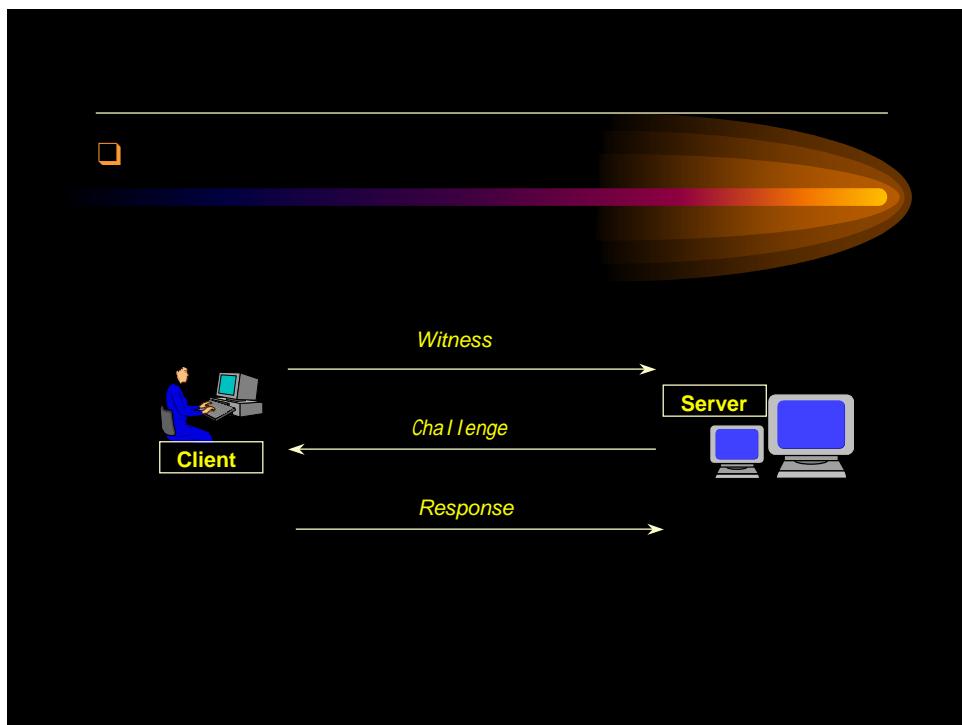
❑ Reflection Attack ()



❑ Zero-Knowledge Proof



❑ Zero-Knowledge Proof of Identity



□ Quadratic Residues modulo a prime

- p is a prime
- If $(v,p) = 1$ and $x^2 \bmod p = v$ has a solution, then v is called a "**quadratic residue**" of modulo p .
 x is a "**square root**" of quadratic residue v .

- e.g.
 $p=7$, $x^2 \bmod p = v$, $(v,p) = 1$
 $1^2 \bmod 7 = 1$, $2^2 \bmod 7 = 4$, $3^2 \bmod 7 = 2$
 $4^2 \bmod 7 = 2$, $5^2 \bmod 7 = 4$, $6^2 \bmod 7 = 1$
 $\{1,2,4\}$ = a set of quadratic residues of modulo 7
 $\{3,5,6\}$ = a set of quadratic non-residues

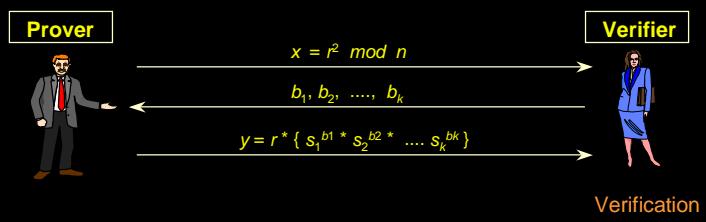
□ Quadratic Residues modulo a prime

- Let $s = (p - 1) / 2$, where $p > 2$ is a prime
- $S_p = \{x^2 \bmod p : 0 < x \leq s\}$
- The elements of S_p are quadratic residues modulo p
- There are s quadratic residues and s quadratic non-residues

- Let $s = (p - 1) / 2$, where $p > 2$ is a prime
- Let v_i be a quadratic residue modulo p
- v_i has exactly two square roots modulo p in \mathbb{Z}_p^*
- one in $(0, s]$ and the other in $(s, p-1]$

□ Feige-Fiat-Shamir Interactive Identification

- a public modulus $n = p * q$
- choose a random v_i such that $x^2 \bmod n = v_i$, where $i = 1, 2, \dots, t$ and $v_i^{-1} \bmod n$ exists
- compute $s_i = \sqrt(v_i^{-1}) \bmod n$



❑ Schnorr

