

## Chapter 7. Introduction to Number Theory

- 7.1 Prime and Relatively Prime Numbers
- 7.2 Modular Arithmetic
- 7.3 Fermat's and Euler's Theorems
- 7.4 Testing for Primality
- 7.5 Euclid's Algorithm
- 7.6 The Chinese Remainder Theorem
- 7.7 Discrete Logarithms

### 7.1 Prime and Relatively Prime (1)

- 1. Divisors

- ✓  $b \mid a$  :  $b (\neq 0)$  divides  $a$  i.e.  $a = mb$  for some  $m$ .
- ✓  $b$  is divisor of  $a$ .

- ✓ Properties

- ✓  $a \mid 1$        $a = \pm 1$
- ✓  $a \mid b$  and  $b \mid a$        $a = \pm b$
- ✓ Any  $b (\neq 0) \mid 0$
- ✓  $b \mid g$  and  $b \mid h$        $b \mid (mg + nh)$  for arbitrary  $m$  and  $n$ .

## 7.1 Prime and Relatively Prime (2)

### 2. Prime Numbers

✓  $p$  is prime number if its only divisors are  $\pm 1$  and  $\pm p$ .

✓  $a$  can be factored in a unique way

$$a = p_1^{\alpha_1} p_2^{\alpha_2} \dots p_t^{\alpha_t}$$

where  $p_i$  : distinct prime,  $\alpha_i > 0$ .

## 7.1 Prime and Relatively Prime (3)

### 3. Relatively Prime Numbers

✓ Greatest Common Divisor

✓  $c = \gcd(a, b)$  if  $c \mid a$ ,  $c \mid b$  and  
 $d \mid a$ ,  $d \mid b$   $d \mid c$ .

✓  $\gcd(a, b) = \gcd(|a|, |b|)$

✓  $a$  and  $b$  are relatively prime if  $\gcd(a, b) = 1$ .

## 7.2 Modular Arithmetic (1)

### 1. Remainder

- ✓ Any integer  $a$  satisfy the following relationship

- ✓ Quotient  $a = qn + r$

- ✓ Remainder  $q = \lfloor a/n \rfloor$

- ✓ The remainder  $r$  is often referred to as a residue.  $0 \leq r < n$

- ✓ Examples

$$11 = 1 \times 7 + 4 \Rightarrow r = 4, \quad 11 \bmod 7 = 4$$

$$-11 = (-2) \times 7 + 3 \Rightarrow r = 3, \quad -11 \bmod 7 = 3$$

## 7.2 Modular Arithmetic (2)

### 2. Congruent modulo $n$

- ✓  $a \equiv b \bmod n$  if  $n \mid (a - b)$

- ✓  $(a \bmod n) = (b \bmod n) \Leftrightarrow a \equiv b \bmod n$

- ✓  $a \equiv b \bmod n \Leftrightarrow b \equiv a \bmod n$

- ✓  $a \equiv b \bmod n$  and  $b \equiv c \bmod n \Rightarrow a \equiv c \bmod n$

- ✓ Examples

$$73 \equiv 4 \bmod 23$$

$$21 \equiv -9 \bmod 10$$

## 7.2 Modular Arithmetic (3)

### 3. Modular arithmetic operations

- ✓  $[(a \bmod n) + (b \bmod n)] \bmod n = (a + b) \bmod n$
- ✓  $[(a \bmod n) - (b \bmod n)] \bmod n = (a - b) \bmod n$
- ✓  $[(a \bmod n) \times (b \bmod n)] \bmod n = (a \times b) \bmod n$
- ✓  $(a + b) \equiv (a + c) \bmod n \Rightarrow b \equiv c \bmod n$
- ✓  $(a \times b) \equiv (a \times c) \bmod n \Rightarrow b \equiv c \bmod n$

if  $\gcd(a, n) = 1$

## 7.2 Modular Arithmetic (4)

### ✓ Examples

- ✓  $[(11 \bmod 8) + (15 \bmod 8)] \bmod 8 = 2 = (11 + 15) \bmod 8$
- ✓  $[(11 \bmod 8) - (15 \bmod 8)] \bmod 8 = 4 = (11 - 15) \bmod 8$
- ✓  $[(11 \bmod 8) \times (15 \bmod 8)] \bmod 8 = 5 = (11 \times 15) \bmod 8$
- ✓  $(5 + 23) \bmod 8 = (5 + 7) \bmod 8 = 23 \bmod 8$
- ✓  $(5 \times 23) \bmod 8 = (5 \times 7) \bmod 8 = 23 \bmod 8 \quad \gcd(5, 8) = 1$
- ✓  $(6 \times 3) \bmod 8 = (6 \times 7) \bmod 8 = 3 \bmod 8 \quad \gcd(6, 8) = 2$

|              | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|--------------|---|---|---|---|---|---|---|---|
| $5i \bmod 8$ | 0 | 5 | 2 | 7 | 4 | 1 | 6 | 3 |
| $6i \bmod 8$ | 0 | 6 | 4 | 2 | 0 | 6 | 4 | 2 |

## 7.2 Modular Arithmetic (5)

### ✓ Additive inverse

✓  $(z + w) \equiv 0 \pmod n \Rightarrow z (= -w)$  is the additive inverse of  $w$ .

### ✓ Addition Modulo 7

| + | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
| 1 | 1 | 2 | 3 | 4 | 5 | 6 | 0 |
| 2 | 2 | 3 | 4 | 5 | 6 | 0 | 1 |
| 3 | 3 | 4 | 5 | 6 | 0 | 1 | 2 |
| 4 | 4 | 5 | 6 | 0 | 1 | 2 | 3 |
| 5 | 5 | 6 | 0 | 1 | 2 | 3 | 4 |
| 6 | 6 | 0 | 1 | 2 | 3 | 4 | 5 |

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## 7.2 Modular Arithmetic (6)

### ✓ Multiplicative inverse

✓  $(z \cdot w) \equiv 1 \pmod n \quad z (= w^{-1})$  is the multiplicative inverse of  $w$ .

### ✓ Multiplicative Modulo 7 $(\gcd(w,n)=1)$

|   | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
| 2 | 0 | 2 | 4 | 6 | 1 | 3 | 5 |
| 3 | 0 | 3 | 6 | 2 | 5 | 1 | 4 |
| 4 | 0 | 4 | 1 | 5 | 2 | 6 | 3 |
| 5 | 0 | 5 | 3 | 1 | 6 | 4 | 2 |
| 6 | 0 | 6 | 5 | 4 | 3 | 2 | 1 |

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## 7.3 Fermat's and Euler's Theorems (1)

### 1. Fermat's Theorem

$$a^{p-1} \equiv 1 \pmod{p}$$

where  $p$  is prime and  $a$  is a positive integer.

#### Example

$$7^2 \equiv 49 \equiv 11 \pmod{19}$$

$$7^4 \equiv 11^2 \equiv 121 \equiv 7 \pmod{19}$$

$$7^8 \equiv 7^2 \equiv 11 \pmod{19}$$

$$7^{16} \equiv 11^2 \equiv 7 \pmod{19}$$

$$7^{18} = 7^{16} \times 7^2 \equiv 7 \times 11 \equiv 1 \pmod{19}$$

## 7.3 Fermat's and Euler's Theorems (2)

### 2. Euler's Totient Function

✓  $\phi(n)$ : the number of positive integers less than  $n$  and relatively prime to  $n$ .

$$\phi(n) = \prod_{i=1}^m (n_i^{\alpha_i} - n_i^{\alpha_i-1})$$

where  $n = \prod_{i=1}^m n_i^{\alpha_i}$ ,  $p_i$  is prime and  $\alpha_i$  is a positive integer.

$$\phi(n) = n \prod_{i=1}^m \left(1 - \frac{1}{p_i}\right)$$

## 7.3 Fermat's and Euler's Theorems (3)

### ✓ Examples

$$✓ \phi(7) = 7 - 1 = 6$$

Above 6 integers are  $\{1, 2, 3, 4, 5, 6\}$ .

$$\begin{aligned} ✓ \phi(21) &= \phi(3 \times 7) \\ &= (3-1)(7-1) \\ &= 12 \end{aligned}$$

Above 12 integers are  $\{1, 2, 4, 5, 8, 10, 11, 13, 16, 17, 19, 20\}$ .

## 7.3 Fermat's and Euler's Theorems (4)

### ● 3. Euler's Theorem

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

where  $\gcd(a, n) = 1$ .

### ✓ A useful alternative form

$$a^{\phi(n)+1} \equiv a \pmod{n}$$

where  $\gcd(a, n) = 1$  and  $\gcd(a, n) \neq 1$ .

## 7.4 Testing for Primality (1)

### 1. Miller -Rabin Test

✓ By Fermat's theorem,

$$0 \equiv a^{n-1} - 1$$

$$\equiv a^{m2^k} - 1$$

where  $m$  is odd.

$$\equiv (a^{m2^{k-1}} + 1)(a^{m2^{k-1}} - 1)$$

$$\equiv (a^{m2^{k-1}} + 1)(a^{m2^{k-2}} + 1)(a^{m2^{k-2}} - 1)$$

$$\equiv (a^{m2^{k-1}} + 1)(a^{m2^{k-2}} + 1) \cdots (a^m + 1)(a^m - 1) \quad (1)$$

## 7.4 Testing for Primality (2)

✓ If  $n$  is prime,  $a$  satisfies (1).

✓ If  $a$  satisfies (1),  $n$  is pseudoprime.

✓ If  $n$  is an odd composite integer, then at most 1/4 of all the number  $a$ ,  $1 \leq a \leq n-1$ , satisfy (1).

✓ If (1) is performed  $s$  times, the probability that  $n$  is prime is at least  $1-(1/4)^s$ .

✓ Usually,  $s = 50$ .

## 7.4 Testing for Primality (3)

### ✓ Algorithm

MillerRabin( $a, n$ )

1. write  $n-1 = m2^k$ , where  $m$  is odd.
2. choose a random integer  $a$ ,  $1 \leq a \leq n-1$
3. compute  $b = a^m \bmod n$
4. if  $b \equiv 1 \pmod{n}$  then  
    answer “ $n$  is pseudoprime” and quit
5. for  $i = 0$  to  $k-1$  do
6.     if  $b \equiv -1 \pmod{n}$  then  
        answer “ $n$  is pseudoprime” and quit
- else  
         $b = b^2 \bmod n$
7. answer “ $n$  is composite”

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## 7.5 Euclid's Algorithm (1)

### ● 1. Euclid's Algorithm

#### ✓ Finding the Greatest Common Divisor

$$a = qb + r \quad 0 \leq r < b$$

$$\gcd(a, b) \mid a \ \& \ \gcd(a, b) \mid b \Rightarrow \gcd(a, b) \mid r$$

$$\therefore \gcd(a, b) = \gcd(b, r)$$

#### ✓ More generally

$$\begin{array}{lll} r_0 = q_1 r_1 + r_2 & 0 < r_2 < r_1 & \gcd(r_0, r_1) \\ r_1 = q_2 r_2 + r_3 & 0 < r_3 < r_2 & = \gcd(r_1, r_2) \\ \vdots & \vdots & \vdots \\ r_{m-2} = q_{m-1} r_{m-1} + r_m & 0 < r_m < r_{m-1} & = \gcd(r_{m-1}, r_m) \\ r_{m-1} = q_m r_m & & = r_m \end{array}$$

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## 7.5 Euclid's Algorithm (2)

### ✓ Algorithm

Euclid( $d, f$ )

1.  $X \leftarrow f; Y \leftarrow d$
2. if  $Y = 0$  return  $X = \gcd(d, f)$
3.  $R = X \bmod Y$
4.  $X \leftarrow Y$
5.  $Y \leftarrow R$
6. goto 2

### ✓ Examples

$\gcd(55, 22) = \gcd(22, 11) = 11$

$\gcd(11, 10) = \gcd(10, 1) = 1$

## 7.5 Euclid's Algorithm (3)

### • 2. Extended Euclid's Algorithm

#### ✓ Finding the Multiplicative Inverse

- ✓ Applying Euclid's algorithm to  $X_3$  and  $Y_3$  holding with  
 $fT_1 + dT_2 = T_3 \quad fX_1 + dX_2 = X_3 \quad fY_1 + dY_2 = Y_3.$

- ✓ If  $\gcd(d, f) = 1$ ,  $Y_3$  of the final step is 0 and  $Y_3$  of the preceding step is 1.

- ✓ Therefore,

$$fY_1 + dY_2 = 1 \Rightarrow dY_2 \equiv 1 \pmod{f}$$

$Y_2$  is the multiplicative inverse of  $d$ , modulo  $f$ .

## 7.5 Euclid's Algorithm (4)

### ✓ Algorithm

ExtendedEuclid( $d, f$ )

1.  $(X1, X2, X3) \leftarrow (1, 0, f); (Y1, Y2, Y3) \leftarrow (0, 1, d)$
2. if  $Y3 = 0$  return  $X3 = \gcd(d, f)$ ; no inverse
3. if  $Y3 = 1$  return  $Y3 = \gcd(d, f); Y2 = d^{-1} \bmod f$
4.  $Q = \lfloor X3/Y3 \rfloor$
5.  $(T1, T2, T3) \leftarrow (X1 - QY1, X2 - QY2, X3 - QY3)$
6.  $(X1, X2, X3) \leftarrow (Y1, Y2, Y3)$
7.  $(Y1, Y2, Y3) \leftarrow (T1, T2, T3)$
8. goto 2

## 7.5 Euclid's Algorithm (5)

✓ Example :  $550^{-1} \bmod 1769 = 550$

| Q | X1  | X2   | X3   | Y1   | Y2   | Y3  |
|---|-----|------|------|------|------|-----|
| - | 1   | 0    | 1769 | 0    | 1    | 550 |
| 3 | 0   | 1    | 550  | 1    | -3   | 119 |
| 4 | 1   | -3   | 119  | -4   | 13   | 74  |
| 1 | -4  | 13   | 74   | 5    | -16  | 45  |
| 1 | 5   | -16  | 45   | -9   | 29   | 29  |
| 1 | -9  | 29   | 29   | 14   | -45  | 16  |
| 1 | 14  | -45  | 16   | -23  | 74   | 13  |
| 1 | -23 | 74   | 13   | 37   | -119 | 3   |
| 4 | 37  | -119 | 3    | -171 | 550  | 1   |

## 7.6 The Chinese Remainder Theorem (1)

### 1. CRT

✓  $M = \prod_{i=1}^k m_i$  where  $m_i$  are pairwise relatively prime.

$A \in Z_M \Rightarrow A \leftrightarrow (a_1, a_2, \dots, a_k)$  where  $a_i = A \bmod m_i$ .

✓  $Z_M \xrightarrow{1-to-1} Z_{m_1} \times Z_{m_2} \times \dots \times Z_{m_k} \quad 1 \leq i \leq k$

✓  $A \rightarrow (a_1, a_2, \dots, a_k)$  is obviously unique.

## 7.6 The Chinese Remainder Theorem (2)

✓  $(a_1, a_2, \dots, a_k) \rightarrow A$

$M_i = M / m_i \quad \text{for } 1 \leq i \leq k$

$c_i = M_i \times (M_i^{-1} \bmod m_i)$

$A = (\sum_{i=1}^k a_i c_i) \bmod M$

$\Rightarrow$  satisfy  $a_i = A \bmod m_i$ , because  $c_i \bmod m_i = 1$ .

✓ Operations of  $Z_M \xleftrightarrow{\text{equivalent}} Z_1 \times Z_2 \times \dots \times Z_k$

## 7.6 The Chinese Remainder Theorem (3)

### ✓ Example

$$\begin{array}{ll}
 \checkmark 1813 & = 37 \times 49 \\
 973 \bmod 1813 & (973 \bmod 37, 973 \bmod 49) \\
 & = (11 \bmod 37, 42 \bmod 49) \\
 + & + \\
 678 \bmod 1813 & (678 \bmod 37, 678 \bmod 49) \\
 & = (12 \bmod 37, 41 \bmod 49) \\
 \Downarrow & \Downarrow \\
 (973 + 678) \bmod 1813 & (11 + 12 \bmod 37, 42 + 41 \bmod 49) \\
 = 1651 \bmod 1813 & = (23 \bmod 37, 34 \bmod 49) \\
 A = a_1 M_1 (M_1^{-1} \bmod m_1) + a_2 M_2 (M_2^{-1} \bmod m_2) & \\
 = [(23)(49)(34) + (34)(37)(4)] \bmod 1813 = 1651 & 
 \end{array}$$

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## 7.7 Discrete Logarithms (1)

### • 1. The order of $a \pmod n$

- ✓ The least positive exponent  $m$  satisfying equation

$$a^m \equiv 1 \pmod n$$

### • 2. A primitive root of $n$

- ✓ If  $m = \phi(n)$ ,  $a$  is a primitive root of  $n$ .

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## 7.7 Discrete Logarithms (2)

### 3. Powers of Integers, Modulo 19

| $a$ | $a^2$ | $a^3$ | $a^4$ | $a^5$ | $a^6$ | $a^7$ | $a^8$ | $a^9$ | $a^{10}$ | $a^{11}$ | $a^{12}$ | $a^{13}$ | $a^{14}$ | $a^{15}$ | $a^{16}$ | $a^{17}$ | $a^{18}$ | order | pri. |
|-----|-------|-------|-------|-------|-------|-------|-------|-------|----------|----------|----------|----------|----------|----------|----------|----------|----------|-------|------|
| 1   | 1     | 1     | 1     | 1     | 1     | 1     | 1     | 1     | 1        | 1        | 1        | 1        | 1        | 1        | 1        | 1        | 1        | 1     | *    |
| 2   | 4     | 8     | 16    | 13    | 7     | 14    | 9     | 18    | 17       | 15       | 11       | 3        | 6        | 12       | 5        | 13       | 1        | 18    | *    |
| 3   | 9     | 8     | 5     | 15    | 7     | 2     | 6     | 18    | 16       | 10       | 11       | 14       | 4        | 12       | 17       | 13       | 1        | 18    | *    |
| 4   | 16    | 7     | 9     | 17    | 11    | 6     | 5     | 1     | 4        | 16       | 7        | 9        | 17       | 11       | 6        | 5        | 1        | 9     |      |
| 5   | 6     | 11    | 17    | 9     | 7     | 16    | 4     | 1     | 5        | 6        | 11       | 17       | 9        | 7        | 16       | 4        | 1        | 9     |      |
| 6   | 17    | 7     | 4     | 5     | 11    | 9     | 16    | 1     | 6        | 17       | 7        | 4        | 5        | 11       | 9        | 16       | 1        | 9     |      |
| 7   | 11    | 1     | 7     | 11    | 1     | 7     | 11    | 1     | 7        | 11       | 1        | 7        | 11       | 1        | 7        | 11       | 1        | 3     |      |
| 8   | 7     | 18    | 11    | 12    | 1     | 8     | 7     | 18    | 11       | 12       | 1        | 8        | 7        | 18       | 11       | 12       | 1        | 6     |      |
| 9   | 5     | 7     | 6     | 16    | 11    | 4     | 17    | 1     | 9        | 5        | 7        | 6        | 16       | 11       | 4        | 17       | 1        | 9     |      |
| 10  | 5     | 12    | 6     | 3     | 11    | 15    | 17    | 18    | 9        | 14       | 7        | 13       | 16       | 8        | 4        | 2        | 1        | 18    | *    |
| 11  | 7     | 1     | 11    | 7     | 1     | 11    | 7     | 1     | 11       | 7        | 1        | 11       | 7        | 1        | 11       | 7        | 1        | 3     |      |
| 12  | 11    | 18    | 7     | 8     | 1     | 12    | 11    | 18    | 7        | 8        | 1        | 12       | 11       | 18       | 7        | 8        | 1        | 6     |      |
| 13  | 17    | 12    | 4     | 14    | 11    | 10    | 16    | 18    | 6        | 2        | 7        | 15       | 5        | 8        | 9        | 3        | 1        | 18    | *    |
| 14  | 6     | 8     | 17    | 10    | 7     | 3     | 4     | 18    | 5        | 13       | 11       | 2        | 9        | 12       | 16       | 15       | 1        | 18    | *    |
| 15  | 16    | 12    | 9     | 2     | 11    | 13    | 5     | 18    | 4        | 3        | 7        | 10       | 17       | 8        | 6        | 14       | 1        | 18    | *    |
| 16  | 9     | 11    | 5     | 4     | 7     | 17    | 6     | 1     | 16       | 9        | 11       | 5        | 4        | 7        | 17       | 6        | 1        | 9     |      |
| 17  | 4     | 11    | 16    | 6     | 7     | 5     | 9     | 1     | 17       | 4        | 11       | 16       | 6        | 7        | 5        | 9        | 1        | 9     |      |
| 18  | 1     | 18    | 1     | 18    | 1     | 18    | 1     | 18    | 1        | 18       | 1        | 18       | 1        | 18       | 1        | 18       | 1        | 2     |      |

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